

NAME

gitformat-index - Git index format

SYNOPSIS

\$GIT_DIR/index

DESCRIPTION

Git index format

THE GIT INDEX FILE HAS THE FOLLOWING FORMAT

All binary numbers are in network byte order.

In a repository using the traditional SHA-1, checksums and object IDs (object names) mentioned below are all computed using SHA-1. Similarly, in SHA-256 repositories, these values are computed using SHA-256.

Version 2 is described here unless stated otherwise.

⊕

12-byte header consisting of

4-byte signature:

The signature is { 'D', 'I', 'R', 'C' } (stands for "dircache")

4-byte version number:

The current supported versions are 2, 3 and 4.

32-bit number of index entries.

⊕

number of sorted index entries (see below).

⊕

Extensions are identified by signature. Optional extensions can be ignored if Git does not understand them.

4-byte extension signature. If the first byte is 'A'..'Z' the extension is optional and can be ignored.

32-bit size of the extension

Extension data

⊕

checksum over the content of the index file before this checksum.

INDEX ENTRY

Index entries are sorted in ascending order on the name field, interpreted as a string of unsigned bytes (i.e. memcmp() order, no localization, no special casing of directory separator '/'). Entries with the same name are sorted by their stage field.

An index entry typically represents a file. However, if sparse-checkout is enabled in cone mode ('core.sparseCheckoutCone' is enabled) and the 'extensions.sparseIndex' extension is enabled, then the index may contain entries for directories outside of the sparse-checkout definition. These entries have mode '040000', include the 'SKIP_WORKTREE' bit, and the path ends in a directory separator.

32-bit ctime seconds, the last time a file's metadata changed
this is stat(2) data

32-bit ctime nanosecond fractions
this is stat(2) data

32-bit mtime seconds, the last time a file's data changed
this is stat(2) data

32-bit mtime nanosecond fractions
this is stat(2) data

32-bit dev
this is stat(2) data

32-bit ino
this is stat(2) data

32-bit mode, split into (high to low bits)

16-bit unused, must be zero

4-bit object type

valid values in binary are 1000 (regular file), 1010 (symbolic link)
and 1110 (gitlink)

3-bit unused, must be zero

9-bit unix permission. Only 0755 and 0644 are valid for regular files.
Symbolic links and gitlinks have value 0 in this field.

32-bit uid

this is stat(2) data

32-bit gid

this is stat(2) data

32-bit file size

This is the on-disk size from stat(2), truncated to 32-bit.

Object name for the represented object

A 16-bit 'flags' field split into (high to low bits)

1-bit assume-valid flag

1-bit extended flag (must be zero in version 2)

2-bit stage (during merge)

12-bit name length if the length is less than 0xFFFF; otherwise 0xFFFF
is stored in this field.

(Version 3 or later) A 16-bit field, only applicable if the
"extended flag" above is 1, split into (high to low bits).

1-bit reserved for future

1-bit skip-worktree flag (used by sparse checkout)

1-bit intent-to-add flag (used by "git add -N")

13-bit unused, must be zero

Entry path name (variable length) relative to top level directory (without leading slash). '/' is used as path separator. The special path components ".", ".." and ".git" (without quotes) are disallowed. Trailing slash is also disallowed.

The exact encoding is undefined, but the '.' and '/' characters are encoded in 7-bit ASCII and the encoding cannot contain a NUL byte (iow, this is a UNIX pathname).

(Version 4) In version 4, the entry path name is prefix-compressed relative to the path name for the previous entry (the very first entry is encoded as if the path name for the previous entry is an empty string). At the beginning of an entry, an integer N in the variable width encoding (the same encoding as the offset is encoded for OFS_DELTA pack entries; see `linkgit:gitformat-pack[5]`) is stored, followed by a NUL-terminated string S. Removing N bytes from the end of the path name for the previous entry, and replacing it with the string S yields the path name for this entry.

1-8 nul bytes as necessary to pad the entry to a multiple of eight bytes while keeping the name NUL-terminated.

(Version 4) In version 4, the padding after the pathname does not exist.

Interpretation of index entries in split index mode is completely different. See below for details.

EXTENSIONS

Cache tree

Since the index does not record entries for directories, the cache entries cannot describe tree objects that already exist in the object database for regions of the index that are unchanged from an existing commit. The cache tree extension stores a recursive tree structure that describes the trees that already exist and completely match sections of the cache entries. This speeds up tree object generation from the index for a new commit by only computing the trees that are "new" to that commit. It also assists when comparing the index to another tree, such

as 'HEAD^{tree}', since sections of the index can be skipped when a tree comparison demonstrates equality.

The recursive tree structure uses nodes that store a number of cache entries, a list of subnodes, and an object ID (OID). The OID references the existing tree for that node, if it is known to exist. The subnodes correspond to subdirectories that themselves have cache tree nodes. The number of cache entries corresponds to the number of cache entries in the index that describe paths within that tree's directory.

The extension tracks the full directory structure in the cache tree extension, but this is generally smaller than the full cache entry list.

When a path is updated in index, Git invalidates all nodes of the recursive cache tree corresponding to the parent directories of that path. We store these tree nodes as being "invalid" by using "-1" as the number of cache entries. Invalid nodes still store a span of index entries, allowing Git to focus its efforts when reconstructing a full cache tree.

The signature for this extension is { 'T', 'R', 'E', 'E' }.

A series of entries fill the entire extension; each of which consists of:

⊕

path component (relative to its parent directory);

⊕

decimal number of entries in the index that is covered by the tree this entry represents (entry_count);

⊕

space (ASCII 32);

⊕

decimal number that represents the number of subtrees this tree has;

⊕

newline (ASCII 10); and

⊕

name for the object that would result from writing this span of index as a tree.

An entry can be in an invalidated state and is represented by having a negative number in the `entry_count` field. In this case, there is no object name and the next entry starts immediately after the newline. When writing an invalid entry, -1 should always be used as `entry_count`.

The entries are written out in the top-down, depth-first order. The first entry represents the root level of the repository, followed by the first subtree--let's call this A--of the root level (with its name relative to the root level), followed by the first subtree of A (with its name relative to A), and so on. The specified number of subtrees indicates when the current level of the recursive stack is complete.

Resolve undo

A conflict is represented in the index as a set of higher stage entries. When a conflict is resolved (e.g. with "git add path"), these higher stage entries will be removed and a stage-0 entry with proper resolution is added.

When these higher stage entries are removed, they are saved in the resolve undo extension, so that conflicts can be recreated (e.g. with "git checkout -m"), in case users want to redo a conflict resolution from scratch.

The signature for this extension is { 'R', 'E', 'U', 'C' }.

A series of entries fill the entire extension; each of which consists of:

⊕

pathname the entry describes (relative to the root of the repository, i.e. full pathname);

⊕

NUL-terminated ASCII octal numbers, entry mode of entries in stage 1 to 3 (a missing stage is represented by "0" in this field); and

⊕

most three object names of the entry in stages from 1 to 3 (nothing is written for a missing stage).

Split index

In split index mode, the majority of index entries could be stored in a separate file. This extension records the changes to be made on top of that to produce the final index.

The signature for this extension is { 'l', 'i', 'n', 'k' }.

The extension consists of:

⊕

of the shared index file. The shared index file path is \$GIT_DIR/sharedindex.<hash>. If all bits are zero, the index does not require a shared index file.

⊕

ewah-encoded delete bitmap, each bit represents an entry in the shared index. If a bit is set, its corresponding entry in the shared index will be removed from the final index. Note, because a delete operation changes index entry positions, but we do need original positions in replace phase, it's best to just mark entries for removal, then do a mass deletion after replacement.

⊕

ewah-encoded replace bitmap, each bit represents an entry in the shared index. If a bit is set, its corresponding entry in the shared index will be replaced with an entry in this index file. All replaced entries are stored in sorted order in this index. The first "1" bit in the replace bitmap corresponds to the first index entry, the second "1" bit to the second entry and so on. Replaced entries may have empty path names to save space.

The remaining index entries after replaced ones will be added to the final index. These added entries are also sorted by entry name then stage.

UNTRACKED CACHE

Untracked cache saves the untracked file list and necessary data to verify the cache. The signature for this extension is { 'U', 'N', 'T', 'R' }.

The extension starts with

⊕

sequence of NUL-terminated strings, preceded by the size of the sequence in variable width encoding. Each string describes the environment where the cache can be used.

⊕

data of \$GIT_DIR/info/exclude. See "Index entry" section from ctime field until "file size".

⊕

data of core.excludesFile

⊕

dir_flags (see struct dir_struct)

⊕

of \$GIT_DIR/info/exclude. A null hash means the file does not exist.

⊕

of core.excludesFile. A null hash means the file does not exist.

⊕

string of per-dir exclude file name. This usually is ".gitignore".

⊕

number of following directory blocks, variable width encoding. If this number is zero, the extension ends here with a following NUL.

⊕

number of directory blocks in depth-first-search order, each consists of

⊕

number of untracked entries, variable width encoding.

⊕

number of sub-directory blocks, variable width encoding.

⊕

directory name terminated by NUL.

⊕

number of untracked file/dir names terminated by NUL.

The remaining data of each directory block is grouped by type:

⊕

ewah
 bitmap,
 the
 n-th
 bit
 marks
 whether
 the
 n-th
 directory
 has
 valid
 untracked
 cache
 entries.

⊕

ewah bitmap, the n-th bit records "check-only" bit of `read_directory_recursive()` for the n-th directory.

⊕

ewah bitmap, the n-th bit indicates whether hash and stat data is valid for the n-th directory and exists in the next data.

⊕

array of stat data. The n-th data corresponds with the n-th "one" bit in the previous ewah bitmap.

⊕

array of hashes. The n-th hash corresponds with the n-th "one" bit in the previous ewah bitmap.

⊕

NUL.

FILE SYSTEM MONITOR CACHE

The file system monitor cache tracks files for which the `core.fsmonitor` hook has told us about changes. The signature for this extension is `{ 'F', 'S', 'M', 'N' }`.

The extension starts with

⊕

version number: the current supported versions are 1 and 2.

⊕

1) 64-bit time: the extension data reflects all changes through the given time which is stored as the nanoseconds elapsed since midnight, January 1, 1970.

⊕

2) A null terminated string: an opaque token defined by the file system monitor application. The extension data reflects all changes relative to that token.

⊕

bitmap size: the size of the CE_FSMONITOR_VALID bitmap.

⊕

ewah bitmap, the n-th bit indicates whether the n-th index entry is not CE_FSMONITOR_VALID.

END OF INDEX ENTRY

The End of Index Entry (EOIE) is used to locate the end of the variable length index entries and the beginning of the extensions. Code can take advantage of this to quickly locate the index extensions without having to parse through all of the index entries.

Because it must be able to be loaded before the variable length cache entries and other index extensions, this extension must be written last. The signature for this extension is { 'E', 'O', 'I', 'E' }.

The extension consists of:

⊕

offset to the end of the index entries

⊕

over the extension types and their sizes (but not their contents). E.g. if we have "TREE" extension that is N-bytes long, "REUC" extension that is M-bytes long, followed by "EOIE", then the hash would be:

$$\text{Hash}(\text{"TREE"} + \langle \text{binary-representation-of-N} \rangle + \text{"REUC"} + \langle \text{binary-representation-of-M} \rangle)$$

INDEX ENTRY OFFSET TABLE

The Index Entry Offset Table (IEOT) is used to help address the CPU

cost of loading the index by enabling multi-threading the process of converting cache entries from the on-disk format to the in-memory format. The signature for this extension is { 'I', 'E', 'O', 'T' }.

The extension consists of:

⊕

version (currently 1)

⊕

number of index offset entries each consisting of:

⊕

offset from the beginning of the file to the first cache entry in this block of entries.

⊕

count of cache entries in this block

SPARSE DIRECTORY ENTRIES

When using sparse-checkout in cone mode, some entire directories within the index can be summarized by pointing to a tree object instead of the entire expanded list of paths within that tree. An index containing such entries is a "sparse index". Index format versions 4 and less were not implemented with such entries in mind. Thus, for these versions, an index containing sparse directory entries will include this extension with signature { 's', 'd', 'i', 'r' }. Like the split-index extension, tools should avoid interacting with a sparse index unless they understand this extension.

GIT

Part of the **git**(1) suite