NAME

vm_map - virtual address space portion of virtual memory subsystem

SYNOPSIS

#include <sys/param.h>
#include <vm/vm.h>
#include <vm/vm_map.h>

DESCRIPTION

The **vm_map** subsystem is used to manage virtual address spaces. This section describes the main data structures used within the code.

The *struct vm_map* is a generic representation of an address space. This address space may belong to a user process or the kernel. The kernel actually uses several maps, which are maintained as subordinate maps, created using the vm_map_submap(9) function.

struct vm_map {
 struct vm_map_entry header;
 struct sx lock;
 struct mtx system_mtx;
 int nentries;
 vm_size_t size;
 u_int timestamp;
 u_char needs_wakeup;
 u_char system_map;
 vm_flags_t flags;
 vm_map_entry_t root;
 pmap_t pmap;
 int busy;
};

The fields of *struct vm_map* are as follows:

- *header* Head node of a circular, doubly linked list of *struct vm_map_entry* objects. Each object defines a particular region within this map's address space.
- *lock* Used to serialize access to the structure.

system_mtx A mutex which is used if the map is a system map.

nentries	A count of the members in use within the circular map entry list.		
size	Specifies the size of the virtual address space.		
timestamp	Used to determine if the map has changed since its last access.		
needs_wakeup	Indicates if a thread is waiting for an allocation within the map. Used only by system maps.		
system_map	Set to TRUE to indicate that map is a system map; otherwise, it belongs to a user process.		
flags	Map flags, described below.		
root	Root node of a binary search tree used for fast lookup of map entries.		
ртар	Pointer to the underlying physical map with which this virtual map is associated.		
busy	Map busy counter, prevents forks.		
Possible map flags:			
MAP_WIREFUTURE		Wire all future pages in this map.	
MAP_BUSY_WAKEUP There		There are waiters for the map busy status.	
The following flags can be passed to vm_map_find(9) and vm_map_insert(9) to specify the copy-on- write properties of regions within the map:			
MAP_COPY_ON_WRITE The mapping is copy-on-write.			

MAP_COPY_ON_WRITE	The mapping is copy-on-write.
MAP_NOFAULT	The mapping should not generate page faults.
MAP_PREFAULT	The mapping should be prefaulted into physical memory.
MAP_PREFAULT_PARTIAL	The mapping should be partially prefaulted into physical memory.
MAP_DISABLE_SYNCER	Do not periodically flush dirty pages; only flush them when absolutely necessary.

MAP_DISABLE_COREDUMP

Do not include the mapping in a core dump.

MAP_PREFAULT_MADVISE Specify that the request is from a user process calling madvise(2).

MAP_ACC_CHARGED Region is already charged to the requestor by some means.

MAP_ACC_NO_CHARGE Do not charge for allocated region.

The *struct vm_map_entry* is a generic representation of a region. The region managed by each entry is associated with a *union vm_map_object*, described below.

struct vm_map_entry { struct vm_map_entry *prev; struct vm_map_entry *next; struct vm_map_entry *left; struct vm_map_entry *right; vm_offset_t start; vm_offset_t end; vm_offset_t avail_ssize; vm_size_t adj_free; vm_size_t max_free; union vm_map_object object; vm_ooffset_t offset; vm_eflags_t eflags; /* Only in task maps: */ vm_prot_t protection; vm_prot_t max_protection; vm_inherit_t inheritance; int wired_count; vm_pindex_t lastr;

};

The fields of *struct vm_map_entry* are as follows:

prev	Pointer to the previous node in a doubly-linked, circular list.
next	Pointer to the next node in a doubly-linked, circular list.
left	Pointer to the left node in a binary search tree.

right Pointer to the right node in a binary search tree.

- *start* Lower address bound of this entry's region.
- *end* Upper address bound of this entry's region.
- *avail_ssize* If the entry is for a process stack, specifies how much the entry can grow.
- *adj_free* The amount of free, unmapped address space adjacent to and immediately following this map entry.
- *max_free* The maximum amount of contiguous free space in this map entry's subtree.
- *object* Pointer to the *struct vm_map_object* with which this entry is associated.
- offset Offset within the *object* which is mapped from *start* onwards.
- *eflags* Flags applied to this entry, described below.

The following five members are only valid for entries forming part of a user process's address space:

protection Memory protection bits applied to this region.

max_protection Mask for the memory protection bits which may be actually be applied to this region.

- *inheritance* Contains flags which specify how this entry should be treated during fork processing.
- *wired_count* Count of how many times this entry has been wired into physical memory.

lastr Contains the address of the last read which caused a page fault.

The following flags may be applied to each entry, by specifying them as a mask within the *eflags* member:

MAP_ENTRY_NOSYNC	The system should not flush the data associated with this map periodically, but only when it needs to.	
MAP_ENTRY_IS_SUB_MAP	If set, then the <i>object</i> member specifies a subordinate map.	
MAP_ENTRY_COW	Indicate that this is a copy-on-write region.	

VM_MAP(9)	FreeBSD	Kernel Developer's Manual	VM_MAP(9)
MAP_ENTRY_NEEDS_COP	Y	Indicate that a copy-on-write region needs to be	copied.
MAP_ENTRY_NOFAULT		Specifies that accesses within this region should a page fault. If a page fault occurs within this re system will panic.	
MAP_ENTRY_USER_WIRE	D	Indicate that this region was wired on behalf of process.	a user
MAP_ENTRY_BEHAV_NOF	RMAL	The system should use the default paging behav region.	iour for this
MAP_ENTRY_BEHAV_SEQ	UENTIAL	The system should depress the priority of pages preceding each page within this region when fat	-
MAP_ENTRY_BEHAV_RAN	NDOM	Is a hint that pages within this region will be according randomly, and that prefetching is likely not adva	
MAP_ENTRY_IN_TRANSIT	ION	Indicate that wiring or unwiring of an entry is in and that other kernel threads should not attempt fields in the structure.	* •
MAP_ENTRY_NEEDS_WAB	KEUP	Indicate that there are kernel threads waiting for become available.	[•] this region to
MAP_ENTRY_NOCOREDU	MP	The region should not be included in a core dun	ıp.

The *inheritance* member has type *vm_inherit_t*. This governs the inheritance behaviour for a map entry during fork processing. The following values are defined for *vm_inherit_t*:

VM_INHERIT_SHARE	The object associated with the entry should be cloned and shared with the new map. A new <i>struct vm_object</i> will be created if necessary.
VM_INHERIT_COPY	The object associated with the entry should be copied to the new map.
VM_INHERIT_NONE	The entry should not be copied to the new map.
VM_INHERIT_DEFAULT	Specifies the default behaviour, VM_INHERIT_COPY.

The union vm_map_object is used to specify the structure which a struct vm_map_entry is associated

with.

The fields of *union vm_map_object* are as follows:

};

Normally, the *sub_map* member is only used by system maps to indicate that a memory range is managed by a subordinate system map. Within a user process map, each *struct vm_map_entry* is backed by a *struct vm_object*.

SEE ALSO

pmap(9), vm_map_check_protection(9), vm_map_delete(9), vm_map_entry_resize_free(9), vm_map_find(9), vm_map_findspace(9), vm_map_inherit(9), vm_map_init(9), vm_map_insert(9), vm_map_lock(9), vm_map_lookup(9), vm_map_madvise(9), vm_map_max(9), vm_map_min(9), vm_map_pmap(9), vm_map_protect(9), vm_map_remove(9), vm_map_simplify_entry(9), vm_map_stack(9), vm_map_submap(9), vm_map_sync(9), vm_map_wire(9)

AUTHORS

This manual page was written by Bruce M Simpson *<bms@spc.org>*.